

## A POSSIBLE BETTER IMPLEMENTATION OF THE BAILLIE-PSW PRIMALITY TEST

**DONATO DI IORIO**

Department of Economics, Management, Society and Institutions  
University of Molise  
via Francesco De Sanctis 1, (86100) Campobasso  
Italy  
e-mail: donato.diiorio@unimol.it

### **Abstract**

In this manuscript, we show first a more synthetic definition of strong pseudoprimality to base 2 and then a possible better implementation of the Baillie-PSW primality test. In particular, about the implementation of the Baillie-PSW primality test, we show that some operations can be avoided [1, 2].

### **1. Introduction**

The main objective of this manuscript is to show the possibility of implementing the Baillie-PSW primality test more appropriately with reference to strong pseudoprimality to base 2. In particular we first show, by means of the characterization of Fermat's little theorem in the

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congruence classes mod 8, a more concise definition of strong pseudoprimality to base 2 and then a possible better implementation of the Baillie-PSW primality test.

## 2. A Necessary Condition of Primality Deriving from Euler's Criterion and Legendre Symbol

Fermat's little theorem states that:

if  $p$  is a prime number and  $a$  is any integer coprime with  $p$ , then it is:

$$a^{p-1} \equiv 1 \pmod{p}.$$

We can express Fermat's little theorem in another way:

if  $p$ ,  $p > 2$ , is a prime number and  $a$  is any integer coprime with  $p$ , then it is:

$$a^{\frac{p-1}{2}} \equiv 1 \pmod{p} \quad \text{or} \quad a^{\frac{p-1}{2}} \equiv -1 \pmod{p}.$$

By means of Euler's criterion and Legendre symbol we can characterize the above condition with respect to the base  $a = 2$  in the congruence classes  $p \equiv 1 \pmod{8}$ ,  $p \equiv 3 \pmod{8}$ ,  $p \equiv 5 \pmod{8}$  and  $p \equiv 7 \pmod{8}$ . In fact, we have the following proposition.

**Proposition 2.1.** *If  $p$ ,  $p > 2$ , is a prime number,  $p \equiv 1 \pmod{8}$  or  $p \equiv 7 \pmod{8}$ , then it is:  $2^{\frac{p-1}{2}} \equiv 1 \pmod{p}$ ; if  $p$ ,  $p > 2$ , is a prime number,  $p \equiv 3 \pmod{8}$  or  $p \equiv 5 \pmod{8}$ , then it is:  $2^{\frac{p-1}{2}} \equiv -1 \pmod{p}$ .*

**Proof.** If  $p$ ,  $p > 2$ , is prime and if, moreover,  $p \equiv 1 \pmod{8}$ , i.e., if  $p = 8k + 1$ ,  $k \in \mathbb{N}$ ,  $k > 1$ , since:

$$\left(\frac{2}{p}\right) = (-1)^{\frac{p^2-1}{8}} = (-1)^{\frac{(8k+1)^2-1}{8}} = (-1)^{\frac{64k^2+1+16k-1}{8}} = (-1)^{2k(4k+1)} = 1,$$

from Euler's criterion  $a^{\frac{p-1}{2}} \equiv \left(\frac{a}{p}\right) \pmod{p}$  for  $a = 2$  it is:

$2^{\frac{p-1}{2}} \equiv 1 \pmod{p}$ . Similarly in other cases  $p \equiv 3 \pmod{8}$ ,  $p \equiv 5 \pmod{8}$  and  $p \equiv 7 \pmod{8}$ .

Proposition 2.1 can also be expressed in the following way.

**Proposition 2.2.** *If  $p$ ,  $p > 2$ , is a prime number such that  $p-1 = 2^s \cdot t$ ,  $s \in \mathbb{N}$ ,  $s \geq 1$ ,  $t \in \mathbb{N}$ ,  $t$  odd, we have the following:*

a) *if  $p \equiv 1 \pmod{8}$  ( $s \geq 3$ ), then we have:*

$2^t \equiv 1 \pmod{p}$  or  $2^{2^r \cdot t} \equiv -1 \pmod{p}$  for some integer  $r$ ,  $0 \leq r \leq s-2$ ;

b) *if  $p \equiv 3 \pmod{8}$ , then we have:  $2^t \equiv -1 \pmod{p}$ ;*

c) *if  $p \equiv 5 \pmod{8}$ , then we have:  $2^{2t} \equiv -1 \pmod{p}$ ;*

d) *if  $p \equiv 7 \pmod{8}$ , then we have:  $2^t \equiv 1 \pmod{p}$ .*

**Proof.** a) From proposition 2.1:  $\left(2^{\frac{p-1}{4}}\right)^2 - 1 \equiv 0 \pmod{p}$ , (in this

case it is  $p \equiv 1 \pmod{8}$ , i.e.,  $p = 8k+1$ ,  $k \in \mathbb{N}$ ,  $k > 1$ ;  $p-1 = 2^3 k$ ,

$k \in \mathbb{N}$ ,  $k > 1$ ; so  $\frac{p-1}{4}$  is an integer number and it is also  $s \geq 3$ ) we

have:

$$\left(2^{\frac{p-1}{4}} + 1\right) \left(2^{\frac{p-1}{8}} + 1\right) \cdot \dots \cdot \left(2^{\frac{p-1}{2^{s-1}}} + 1\right) \left(2^{\frac{p-1}{2^s}} + 1\right) \left(2^{\frac{p-1}{2^s}} - 1\right)$$

$$\equiv 0 \pmod{p};$$

so, since  $\frac{p-1}{2^s} = t$ , we have:

$$(2^t - 1)(2^t + 1)(2^{2t} + 1) \cdot \dots \cdot (2^{2^{s-3} \cdot t} + 1)(2^{2^{s-2} \cdot t} + 1) \equiv 0 \pmod{p},$$

i.e.:  $2^t \equiv 1 \pmod{p}$  or  $2^{2^r \cdot t} \equiv -1 \pmod{p}$  for some integer  $r$ ,  
 $0 \leq r \leq s-2$ ;

b) since it is  $p-1 = 2(4k+1) = 2t$ ,  $k \in N$ , i.e.,  $t = \frac{p-1}{2}$ , from Proposition 2.1 we have:  $2^t \equiv -1 \pmod{p}$ ;

c) since it is  $p-1 = 2^2(2k+1) = 2^2t$ ,  $k \in N$ , i.e.,  $2t = \frac{p-1}{2}$ , from Proposition 2.1 we have:  $2^{2t} \equiv -1 \pmod{p}$ ;

d) since it is  $p-1 = 2(4k+3) = 2t$ ,  $k \in N$ , i.e.,  $t = \frac{p-1}{2}$ , from Proposition 2.1 we have:  $2^t \equiv 1 \pmod{p}$ .

### 3. The Classical Necessary Condition of Primality Deriving from Fermat's Little Theorem

On prime numbers we have the following proposition.

**Proposition 3.1.** *If  $p$ ,  $p > 2$ , is a prime number such that  $p-1 = 2^s \cdot t$ ,  $s \in N$ ,  $s \geq 1$ ,  $t \in N$ ,  $t$  odd, for each integer  $a$  coprime with  $p$ ,  $1 \leq a < p$ , we have:*

$$a^t \equiv 1 \pmod{p} \text{ or } a^{2^r \cdot t} \equiv -1 \pmod{p} \text{ for some integer } r,$$

$$0 \leq r \leq s-1.$$

**Proof.** From Fermat's little theorem:  $\left(a^{\frac{p-1}{2}}\right)^2 - 1 \equiv 0 \pmod{p}$ , it is:

$$\left(a^{\frac{p-1}{2}} + 1\right) \left(a^{\frac{p-1}{4}} + 1\right) \cdot \dots \cdot \left(a^{\frac{p-1}{2^{s-1}}} + 1\right) \left(a^{\frac{p-1}{2^s}} + 1\right) \left(a^{\frac{p-1}{2^s}} - 1\right) \\ \equiv 0 \pmod{p};$$

so, since  $\frac{p-1}{2^s} = t$ , we have:

$$(a^t - 1)(a^t + 1)(a^{2t} + 1) \cdot \dots \cdot (a^{2^{s-2}t} + 1)(a^{2^{s-1}t} + 1) \equiv 0 \pmod{p},$$

i.e.:  $a^t \equiv 1 \pmod{p}$  or  $a^{2^r t} \equiv -1 \pmod{p}$  for some integer  $r$ ,  
 $0 \leq r \leq s-1$ .

With reference to the congruence classes  $p \equiv 1 \pmod{8}$ ,  
 $p \equiv 3 \pmod{8}$ ,  $p \equiv 5 \pmod{8}$  and  $p \equiv 7 \pmod{8}$ , Proposition 3.1  
 becomes the following.

**Proposition 3.2.** *If  $p$ ,  $p > 2$ , is a prime number such that  $p-1 = 2^s \cdot t$ ,  $s \in \mathbb{N}$ ,  $s \geq 1$ ,  $t \in \mathbb{N}$ ,  $t$  odd, for each integer  $a$  coprime with  $p$ ,  $1 \leq a < p$ , we have:*

a) *if  $p \equiv 1 \pmod{8}$  ( $s \geq 3$ ), then we have:  $a^t \equiv 1 \pmod{p}$  or  $a^{2^r t} \equiv -1 \pmod{p}$  for some integer  $r$ ,  $0 \leq r \leq s-1$ ;*

b) *if  $p \equiv 3 \pmod{8}$ , then we have:  $a^t \equiv 1 \pmod{p}$  or  $a^t \equiv -1 \pmod{p}$ ;*

c) *if  $p \equiv 5 \pmod{8}$ , then we have:  $a^t \equiv 1 \pmod{p}$  or*

$a^t \equiv -1 \pmod{p}$  or  $a^{2^t} \equiv -1 \pmod{p}$ ;

d) if  $p \equiv 7 \pmod{8}$ , then we have:  $a^t \equiv 1 \pmod{p}$  or  $a^t \equiv -1 \pmod{p}$ .

**Proof.** a) As in Proposition 3.1, (since  $p \equiv 1 \pmod{8}$ , we have  $p - 1 = 2^3 \cdot k$ ,  $k \in N$ ,  $k > 1$ , that is,  $s \geq 3$ );

b) in fact, since  $p = 8k + 3$ ,  $k \in N$ , i.e.,  $p - 1 = 2(4k + 1)$ ,  $k \in N$ , in reference to Proposition 3.1 we have  $s = 1$ , that is,  $r = 0$ ;

c) in fact, since  $p = 8k + 5$ ,  $k \in N$ , i.e.,  $p - 1 = 2^2(2k + 1)$ ,  $k \in N$ , in reference to Proposition 3.1 we have  $s = 2$ , that is,  $r = 0$  or  $r = 1$ ;

d) in fact, since  $p = 8k + 7$ ,  $k \in N$ , i.e.,  $p - 1 = 2(4k + 3)$ ,  $k \in N$ , in reference to Proposition 3.1 we have  $s = 1$ , that is,  $r = 0$ .

#### 4. On Strong Pseudoprimality to Base 2

Since there are some odd composite integers  $n$  that verify the conditions of Proposition 3.1, we can define the strong pseudoprimality as follows.

**Definition 4.1.** If  $n$  is an odd composite integer, such that  $n - 1 = 2^s \cdot t$ ,  $s \in N$ ,  $s \geq 1$ ,  $t \in N$ ,  $t$  odd, then  $n$  is a strong pseudoprime to the integer base  $a$  ( $spsp(a)$ ),  $1 \leq a < n$ , coprime with  $n$ , if we have:

$a^t \equiv 1 \pmod{n}$  or  $a^{2^r \cdot t} \equiv -1 \pmod{n}$  for some integer  $r$ ,  $0 \leq r \leq s - 1$  (see [2] and [4]).

Regarding to Proposition 3.2, we have the following definition of

strong pseudoprimality to integer base  $a$ ,  $1 \leq a < n$ , coprime with  $n$ , in reference to congruence classes  $n \equiv 1 \pmod{8}$ ,  $n \equiv 3 \pmod{8}$ ,  $n \equiv 5 \pmod{8}$  and  $n \equiv 7 \pmod{8}$ .

**Definition 4.2.** If  $n$  is an odd composite integer, such that  $n - 1 = 2^s \cdot t$ ,  $s \in \mathbb{N}$ ,  $s \geq 1$ ,  $t \in \mathbb{N}$ ,  $t$  odd, then  $n$  is a strong pseudoprime to integer base  $a$  ( $spsp(a)$ ),  $1 \leq a < n$ , coprime with  $n$ , if we have:

$$n \equiv 1 \pmod{8} \quad (s \geq 3) \text{ and}$$

$a^t \equiv 1 \pmod{n}$  or  $a^{2^r \cdot t} \equiv -1 \pmod{n}$  for some integer  $r$ ,  $0 \leq r \leq s - 1$   
or

$$n \equiv 3 \pmod{8} \quad \text{and} \quad (a^t \equiv 1 \pmod{n} \quad \text{or} \quad a^t \equiv -1 \pmod{n})$$

or

$$n \equiv 5 \pmod{8} \text{ and}$$

$$a^t \equiv 1 \pmod{n} \text{ or } a^t \equiv -1 \pmod{n} \text{ or } a^{2t} \equiv -1 \pmod{n}$$

or

$$n \equiv 7 \pmod{8} \text{ and } (a^t \equiv 1 \pmod{n} \text{ or } a^t \equiv -1 \pmod{n}).$$

Since some odd composite integers  $n$  satisfy the conditions of Proposition 2.2, it is possible to define the strong pseudoprimality to base 2 in a more synthetic way than Definition 4.2 with  $a = 2$  as follows.

**Definition 4.3.** If  $n$  is an odd composite integer, such that  $n - 1 = 2^s \cdot t$ ,  $s \in \mathbb{N}$ ,  $s \geq 1$ ,  $t \in \mathbb{N}$ ,  $t$  odd, then  $n$  is a strong pseudoprime to base 2 ( $spsp(2)$ ) if we have:

$$n \equiv 1 \pmod{8} \quad (s \geq 3) \text{ and}$$

$$2^t \equiv 1 \pmod{n} \text{ or } 2^{2^r \cdot t} \equiv -1 \pmod{n} \text{ for some integer } r, \quad 0 \leq r \leq s - 2$$

or

$$n \equiv 3 \pmod{8} \quad \text{and} \quad 2^t \equiv -1 \pmod{n}$$

or

$$n \equiv 5 \pmod{8} \quad \text{and} \quad 2^{2t} \equiv -1 \pmod{n}$$

or

$$n \equiv 7 \pmod{8} \quad \text{and} \quad 2^t \equiv 1 \pmod{n}.$$

**Remark 4.1.** If  $n$ ,  $n > 2$ , is an odd integer such that  $n - 1 = 2^s \cdot t$ ,  $s \in N$ ,  $s \geq 1$ ,  $t \in N$ ,  $t$  odd, considered the proof of Proposition 2.2, to calculate  $n \equiv 1 \pmod{8}$ ,  $n \equiv 3 \pmod{8}$ ,  $n \equiv 5 \pmod{8}$  and  $n \equiv 7 \pmod{8}$ , it is sufficient to compute only  $s$  and  $t$ . In fact we have:

$$n \equiv 1 \pmod{8} \Leftrightarrow s \geq 3; \quad n \equiv 3 \pmod{8} \Leftrightarrow s = 1 \text{ and } t = 4k + 1, \quad k \in N;$$

$$n \equiv 5 \pmod{8} \Leftrightarrow s = 2; \quad n \equiv 7 \pmod{8} \Leftrightarrow s = 1 \text{ and } t = 4k + 3, \quad k \in N.$$

We can give Definition 4.3, using only the values  $s$  and  $t$ .

**Definition 4.4.** If  $n$  is an odd composite integer, such that  $n - 1 = 2^s \cdot t$ ,  $s \in N$ ,  $s \geq 1$ ,  $t \in N$ ,  $t$  odd, then  $n$  is a strong pseudoprime to base 2 ( $spsp(2)$ ) if we have:

$$s \geq 3 \quad (n \equiv 1 \pmod{8}) \text{ and}$$

$$2^t \equiv 1 \pmod{n} \text{ or } 2^{2^r \cdot t} \equiv -1 \pmod{n} \text{ for some integer } r, \quad 0 \leq r \leq s - 2$$

or

$$s = 1 \text{ and } t \equiv 1 \pmod{4} \quad (n \equiv 3 \pmod{8}) \text{ and } 2^t \equiv -1 \pmod{n}$$

or

$$s = 2 \quad (n \equiv 5 \pmod{8}) \text{ and } 2^{2t} \equiv -1 \pmod{n}$$



or

$$s = 1 \text{ and } t \equiv 3 \pmod{4} \text{ } (n \equiv 7 \pmod{8}) \text{ and } 2^t \equiv 1 \pmod{n}.$$

#### 4.1. Some examples on the application of Proposition 2.2 and Definition 4.4

In this section, we study some odd integers, using Proposition 2.2 and Definition 4.4.

**Example 4.1.** Considering  $n = 220729$ , it is:  $n > 2$ ,  $n - 1 = 220728 = 2^3 \cdot 27591$ ,  $s = 3$ ,  $t = 27591$ ,  $220729 \equiv 1 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{27591} \equiv 1 \pmod{220729},$$

$n = 220729 = 103 \cdot 2143$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.2.** Considering  $n = 280601$ , it is:  $n > 2$ ,  $n - 1 = 280600 = 2^3 \cdot 35075$ ,  $s = 3$ ,  $t = 35075$ ,  $280601 \equiv 1 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{35075} \equiv 251179 \not\equiv 1 \pmod{280601},$$

$$2^t = 2^{35075} \equiv 251179 \not\equiv -1 \pmod{280601},$$

$$2^{2^{s-2} \cdot t} = 2^{2^t} = 2^{2 \cdot 35075} \equiv -1 \pmod{280601},$$

$n = 280601 = 277 \cdot 1013$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.3.** Considering  $n = 396271$ , it is:  $n > 2$ ,  $n - 1 = 396270 = 2 \cdot 198135$ ,  $s = 1$ ,  $t = 198135$ ,  $198135 \equiv 3 \pmod{4}$ ,  $396271 \equiv 7 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{198135} \equiv 282542 \not\equiv 1 \pmod{396271},$$

for the counternominal proposition of Proposition 2.2,  $n = 396271$  is a composite integer. Furthermore,  $n = 396271 = 223 \cdot 1777$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.4.** Considering  $n = 489997$ , it is:  $n > 2$ ,  $n - 1 = 489996 = 2^2 \cdot 122499$ ,  $s = 2$ ,  $t = 122499$ ,  $489997 \equiv 5 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{122499} \equiv 249759 \pmod{489997},$$

$$2^{2t} = 2^{2 \cdot 122499} \equiv -1 \pmod{489997},$$

$n = 489997 = 157 \cdot 3121$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.5.** Considering  $n = 877099$ , it is:  $n > 2$ ,  $n - 1 = 877098 = 2 \cdot 438549$ ,  $s = 1$ ,  $t = 438549$ ,  $438549 \equiv 1 \pmod{4}$ ,  $877099 \equiv 3 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{438549} \equiv -1 \pmod{877099},$$

$n = 877099 = 307 \cdot 2857$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.6.** Considering  $n = 3828001$ , it is:  $n > 2$ ,  $n - 1 = 3828000 = 2^5 \cdot 119625$ ,  $s = 5$ ,  $t = 119625$ ,  $3828001 \equiv 1 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{119625} \equiv 2879722 \not\equiv 1 \pmod{3828001},$$

$$2^t = 2^{119625} \equiv 2879722 \not\equiv -1 \pmod{3828001},$$

$$2^{2t} = 2^{2 \cdot 119625} \equiv 1174932 \not\equiv -1 \pmod{3828001},$$

$$2^{2^2 \cdot t} = 2^{2^2 \cdot 119625} \equiv 1 \not\equiv -1 \pmod{3828001},$$

$$2^{2^{s-2} \cdot t} = 2^{2^3 \cdot t} = 2^{2^3 \cdot 119625} \equiv 1 \not\equiv -1 \pmod{3828001},$$

for the counternominal proposition of Proposition 2.2,  $n = 3828001$  is a composite integer. Furthermore,  $n = 3828001 = 101 \cdot 151 \cdot 251$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.7.** Considering  $n = 1251949$ , it is:  $n > 2$ ,  $n - 1 = 1251948 = 2^2 \cdot 312987$ ,  $s = 2$ ,  $t = 312987$ ,  $1251949 \equiv 5 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{312987} \equiv 755566 \pmod{1251949},$$

$$2^{2t} = 2^{2 \cdot 312987} \equiv -1 \pmod{1251949},$$

$n = 1251949 = 409 \cdot 3061$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.8.** Considering  $n = 3421589$ , it is:  $n > 2$ ,  $n - 1 = 3421588 = 2^2 \cdot 855397$ ,  $s = 2$ ,  $t = 855397$ ,  $3421589 \equiv 5 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{855397} \equiv 2301490 \pmod{3421589},$$

$$2^{2t} = 2^{2 \cdot 855397} \equiv 358459 \not\equiv -1 \pmod{3421589},$$

for the counternominal proposition of Proposition 2.2,  $n = 3421589$  is a composite integer. Furthermore,  $n = 3421589 = 131 \cdot 26119$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.9.** Considering  $n = 29111881$ , it is:  $n > 2$ ,  $n - 1 = 29111880 = 2^3 \cdot 3638985$ ,  $s = 3$ ,  $t = 3638985$ ,  $29111881 \equiv 1 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{3638985} \equiv -1 \pmod{29111881},$$

$n = 29111881 = 211 \cdot 281 \cdot 491$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.10.** Considering  $n = 19384289$ , it is:  $n > 2$ ,  $n - 1 = 19384288 = 2^5 \cdot 605759$ ,  $s = 5$ ,  $t = 605759$ ,  $19384289 \equiv 1 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{605759} \equiv 16784867 \not\equiv 1 \pmod{19384289},$$

$$2^t = 2^{605759} \equiv 16784867 \not\equiv -1 \pmod{19384289},$$

$$2^{2t} = 2^{2 \cdot 605759} \equiv 19274464 \not\equiv -1 \pmod{19384289},$$

$$2^{2^2 \cdot t} = 2^{2^2 \cdot 605759} \equiv 4502867 \not\equiv -1 \pmod{19384289},$$

$$2^{2^{s-2} \cdot t} = 2^{2^3 \cdot t} = 2^{2^3 \cdot 605759} \equiv 1 \not\equiv -1 \pmod{19384289},$$

for the counternominal proposition of Proposition 2.2,  $n = 19384289$  is a composite integer. Furthermore,  $n = 19384289 = 89 \cdot 353 \cdot 617$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.11.** Considering  $n = 15247621$ , it is:  $n > 2$ ,  $n - 1 = 15247620 = 2^2 \cdot 3811905$ ,  $s = 2$ ,  $t = 3811905$ ,  $15247621 \equiv 5 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{3811905} \equiv 9141205 \pmod{15247621},$$

$$2^{2t} = 2^{2 \cdot 3811905} \equiv -1 \pmod{15247621},$$

$n = 15247621 = 61 \cdot 181 \cdot 1381$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.12.** Considering  $n = 612816751$ , it is:  $n > 2$ ,  $n - 1 = 612816750 = 2 \cdot 306408375$ ,  $s = 1$ ,  $t = 306408375$ ,  $306408375 \equiv 3 \pmod{4}$ ,  $612816751 \equiv 7 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{306408375} \equiv 550800674 \not\equiv 1 \pmod{612816751},$$

for the counternominal proposition of Proposition 2.2,  $n = 612816751$  is a composite integer. Furthermore,  $n = 612816751 = 251 \cdot 751 \cdot 3251$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.13.** Considering  $n = 7279379941$ , it is:  $n > 2$ ,  $n - 1 = 7279379940 = 2^2 \cdot 1819844985$ ,  $s = 2$ ,  $t = 1819844985$ ,  $7279379941 \equiv 5 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{1819844985} \equiv 852187010 \pmod{7279379941},$$

$$2^{2t} = 2^{2 \cdot 1819844985} \equiv 4443277458 \not\equiv -1 \pmod{7279379941},$$

for the counternominal proposition of Proposition 2.2,  $n = 7279379941$  is a composite integer. Furthermore,  $n = 7279379941 = 211 \cdot 3571 \cdot 9661$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.14.** Considering  $n = 11239359601$ , it is:  $n > 2$ ,  $n - 1 = 11239359600 = 2^4 \cdot 702459975$ ,  $s = 4$ ,  $t = 702459975$ ,  $11239359601 \equiv 1 \pmod{8}$ ; moreover, since it is:

$$2^t = 2^{702459975} \equiv 6448799664 \not\equiv 1 \pmod{11239359601},$$

$$2^t = 2^{702459975} \equiv 6448799664 \not\equiv -1 \pmod{11239359601},$$

$$2^{2t} = 2^{2 \cdot 702459975} \equiv 5391256728 \not\equiv -1 \pmod{11239359601},$$

$$2^{2^{s-2} \cdot t} = 2^{2^2 \cdot t} = 2^{2^2 \cdot 702459975} \equiv 1 \not\equiv -1 \pmod{11239359601},$$

for the counternominal proposition of Proposition 2.2,  $n = 11239359601$  is a composite integer. Furthermore  $n = 11239359601 = 281 \cdot 4201 \cdot 9521$  is not a  $spsp(2)$  (Def. 4.4).

**Example 4.15.** Considering  $n = 83828294551$ , it is:  $n > 2$ ,  $n - 1 = 83828294550 = 2 \cdot 41914147275$ ,  $s = 1$ ,  $t = 41914147275$ ,  $41914147275 \equiv 3 \pmod{4}$ ,  $83828294551 \equiv 7 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{41914147275} \equiv 1 \pmod{83828294551},$$

$n = 83828294551 = 1231 \cdot 6151 \cdot 11071$  is a  $spsp(2)$  (Def. 4.4).

**Example 4.16.** Considering  $n = 3215031751$ , it is:  $n > 2$ ,  $n - 1 = 3215031750 = 2 \cdot 1607515875$ ,  $s = 1$ ,  $t = 1607515875$ ,  $1607515875 \equiv 3 \pmod{4}$ ,  $3215031751 \equiv 7 \pmod{8}$ ; moreover, since we have:

$$2^t = 2^{1607515875} \equiv 1 \pmod{3215031751},$$

$n = 3215031751 = 151 \cdot 751 \cdot 28351$  is a  $spsp(2)$  (Def. 4.4).

## 5. A Possible Better Implementation of the Baillie-PSW Primality Test

The Baillie-PSW primality test (Pomerance 1984) is a probabilistic algorithm to study the primality of odd integers  $n$ ,  $n > 2$ , which consists of the following steps (see [1] and [3]).

### Algorithm 1:

a) A strong pseudoprimal test to base 2 is performed (Definition 4.1 with  $a = 2$ ); if the test is not verified, for the counternominal proposition of Proposition 3.1 with  $a = 2$ ,  $n$  is a composite integer and the Algorithm 1 stops, otherwise, if it is verified, since  $n$  can be a prime number or a strong pseudoprime to base 2, according to Definition 4.1, you go on to next step;

b) In the sequence  $5, -7, 9, -11, \dots$  the first number  $D$  for which

the symbol of Jacobi  $\left(\frac{D}{n}\right) = -1$  is found; then a Lucas pseudoprimality test with discriminant  $D$  on  $n$  is performed. If the test is not verified  $n$  is a composite integer, otherwise,  $n$  is most likely prime.

To improve the above implementation of the Baillie-PSW primality test, with reference to the version of Pomerance (1984) (see [3]), it is possible to apply initially, instead of the strong pseudoprimality test to the base 2, according to the Definition 4.1 with  $a = 2$ , the strong pseudoprimality test to base 2, according to Definition 4.4. So in detail if  $n$ ,  $n > 2$ , is an odd integer such that  $n - 1 = 2^s \cdot t$ ,  $s \in \mathbb{N}$ ,  $s \geq 1$ ,  $t \in \mathbb{N}$ ,  $t$  odd, Algorithm 1 becomes the following, which is more synthetic.

**Algorithm 2:**

$\alpha$ ) If  $s \geq 3$  ( $n \equiv 1 \pmod{8}$ ) you check if it is:

$2^t \equiv 1 \pmod{n}$  or  $2^{2^{r \cdot t}} \equiv -1 \pmod{n}$  for some integer  $r$ ,  $0 \leq r \leq s - 2$  (5.1); if condition (5.1) is not verified, for the counternominal proposition of Proposition 2.2,  $n$  is a composite integer and the Algorithm 2 stops, otherwise, if it is verified, since  $n$  can be a prime number or a strong pseudoprime to base  $p_1 = 2$ , according to Definition 4.4, you go on to next step;

$\alpha_1$ ) You apply Step b) of the Algorithm 1;

$\beta$ ) If  $s = 1$  and  $t \equiv 1 \pmod{4}$  ( $n \equiv 3 \pmod{8}$ ), you check if it is:

$$2^t \equiv -1 \pmod{n}; \quad (5.2)$$

if condition (5.2) is not verified, for the counternominal proposition of Proposition 2.2,  $n$  is a composite integer and the Algorithm 2 stops, otherwise, if it is verified, since  $n$  can be a prime number or a strong

pseudoprime to base  $p_1 = 2$ , according to Definition 4.4, you go on to next step;

$\beta_1$ ) You apply Step b) of the Algorithm 1;

$\gamma$ ) If  $s = 2$  ( $n \equiv 5 \pmod{8}$ ) you check if it is:  $2^{2t} \equiv -1 \pmod{n}$ ; (5.3)

if condition (5.3) is not verified, for the counternominal proposition of Proposition 2.2,  $n$  is a composite integer and the Algorithm 2 stops, otherwise, if it is verified, since  $n$  can be a prime number or a strong pseudoprime to base  $p_1 = 2$ , according to Definition 4.4, you go on to next step;

$\gamma_1$ ) You apply Step b) of the Algorithm 1;

$\delta$ ) If  $s = 1$  and  $t \equiv 3 \pmod{4}$  ( $n \equiv 7 \pmod{8}$ ), you check if it is:

$$2^t \equiv 1 \pmod{n}; \quad (5.4)$$

if condition (5.4) is not verified, for the counternominal proposition of Proposition 2.2,  $n$  is a composite integer and the Algorithm 2 stops, otherwise, if it is verified, since  $n$  can be a prime number or a strong pseudoprime to base  $p_1 = 2$  according to Definition 4.4, you go on to next step;

$\delta_1$ ) You apply Step b) of the Algorithm 1.

## 6. Conclusions

If  $n$ ,  $n > 2$ , is an odd integer such that  $n - 1 = 2^s \cdot t$ ,  $s \in N$ ,  $s \geq 1$ ,  $t \in N$ ,  $t$  odd, considering Algorithm 1 (Pomerance 1984) and Algorithm 2, related to the Baillie-PSW primality test, comparing Definition 4.2 with  $a = 2$  (see Definition 4.1 with  $a = 2$ ) and Definition 4.4, we can state that:



a) if  $s \geq 3$  ( $n \equiv 1 \pmod{8}$ ) it is not necessary to check if it is:

$$2^t \equiv 1 \pmod{n} \text{ or } 2^{2^r \cdot t} \equiv -1 \pmod{n} \text{ for some integer } r, \ 0 \leq r \leq s-1,$$

since it is sufficient to check only:

$$2^t \equiv 1 \pmod{n} \text{ or } 2^{2^{s-1} \cdot t} \equiv -1 \pmod{n} \text{ for some integer } r, \ 0 \leq r \leq s-2;$$

so it is not necessary to check if it is:  $2^{2^{s-1} \cdot t} \equiv -1 \pmod{n}$ ;

b) if  $s = 1$  and  $t \equiv 1 \pmod{4}$  ( $n \equiv 3 \pmod{8}$ ), it is not necessary to check if it is:

$$2^t \equiv 1 \pmod{n} \quad \text{or} \quad 2^t \equiv -1 \pmod{n},$$

since it is sufficient to check only:  $2^t \equiv -1 \pmod{n}$ ;

c) if  $s = 2$  ( $n \equiv 5 \pmod{8}$ ), it is not necessary to check if it is:

$$2^t \equiv 1 \pmod{n} \quad \text{or} \quad 2^t \equiv -1 \pmod{n} \quad \text{or} \quad 2^{2t} \equiv -1 \pmod{n},$$

since it is sufficient to check only:  $2^{2t} \equiv -1 \pmod{n}$ ;

d) if  $s = 1$  and  $t \equiv 3 \pmod{4}$  ( $n \equiv 7 \pmod{8}$ ), it is not necessary to check if it is:

$$2^t \equiv 1 \pmod{n} \quad \text{or} \quad 2^t \equiv -1 \pmod{n},$$

since it is sufficient to check only:  $2^t \equiv 1 \pmod{n}$ .

Therefore, some unnecessary checks can be avoided in the implementation of the Baillie-PSW primality test.

**References**

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